

# Equal opportunities in school choice

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# Motivation

- Equality of opportunity (EOp) - (*Roemer 1998, Fleurbaey, 2008*)
- Education key for EOp - (*Roemer and Unveren, 2017; Corak, 2013*)
- Typical policy recommendation: investment to improve quality.
- What about the students' allocation? (Complementary policy)
- Public school choice is a many-to-one matching market: no price

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# Example: Student proposing Deferred Acceptance

New York City assigns students to High school using Student proposing Deferred Acceptance (Gale and Shapley, 1962)

- $i_1, i_2, i_3 : s_1 P s_2$
- $i_4 : s_2 P s_1$
  
- $s_1 : i_4 \succ i_3 \succ i_2 \succ i_1; \quad q = 1$
- $s_2 : i_1 \succ i_2 \succ i_3 \succ i_4; \quad q = 3$
  
- Round 1:
  - ▶  $s_1 \leftarrow (j_1, j_2, i_3)$
  - ▶  $s_2 \leftarrow (i_4)$
- Round 2:
  - ▶  $s_1 \leftarrow (i_3)$
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# Affirmative Action and Match Quality

- Distributional concern in the matching theory literature: affirmative action.
- Affirmative action  $\neq$  EOp
  - ▶ Who is losing? (Students for Fair Admissions, Inc. v. President and Fellows of Harvard College.)
  - ▶ Weakening stability
- Preferences  $\neq$  Quality (*Abdulkadiroglu et al. 2021*)

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# Framework

# The school choice setting

- $\mathcal{I} = \langle I, S, P, \succ, q, T, \triangleleft, U \rangle$
- $I = \{i_1, i_2, \dots, i_{|I|}\}$  students
- $S = \{s_1, s_2, \dots, s_{|S|}\}$  schools
- $P = (P_{i_1}, \dots, P_{i_{|I|}})$  students' preferences profile (strict and complete linear orders) over  $S$
- $\succ = (\succ_{s_1}, \dots, \succ_{s_{|S|}})$  schools' priority
- $q = (q_{s_1}, \dots, q_{s_{|S|}})$  quota profile of schools
- $T = \{t_1, t_2, \dots, t_{|T|}\}$  set of types (or type partition)
- $\triangleleft$  complete pre-order of types s.t.  $t \triangleleft t'$  if students of type  $t$  are less advantaged (or have stronger needs) than those of  $t'$ .
- Match quality

$$\begin{bmatrix} U(i_1, s_1) & \dots & U(i_1, s_{|S|}) \\ \dots & \dots & \dots \\ U(i_{|I|}, s_1) & \dots & U(i_{|I|}, s_{|S|}) \end{bmatrix}$$

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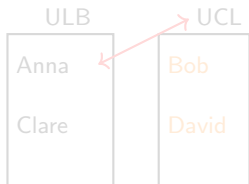
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# Matchings

- An assignment is a function  $\mu : I \cup S \rightarrow 2^{I \cup S}$  such that:
  - ▶ for all  $i \in I$ ,  $\mu(i) \in S \setminus \emptyset$ ;
  - ▶ for all  $s \in S$ ,  $|\mu(s)| \leq q_s$  and  $\mu(s) \subseteq I$ ;
  - ▶  $\mu(i) = s$  if and only if  $i \in \mu(s)$
- **Stability:** there exist no student-school pair preferring each other to their current match.

No blocking pairs:

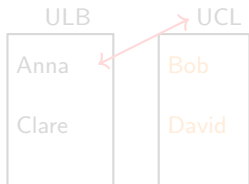


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- BUT enumerating stable assignments is NP-complete (Irving and Leather, 1986)

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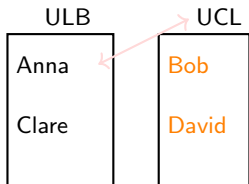


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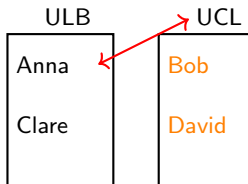


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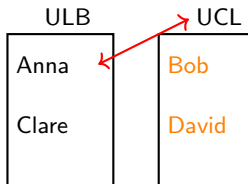


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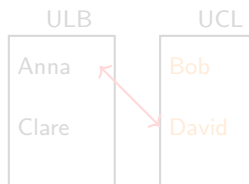
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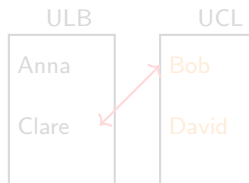
# Desirable requirements

- Efficiency: exhaust Paretian improvements of MQ



$$U(\text{Anna}, \text{ULB}) = U(\text{Anna}, \text{UCL})$$
$$U(\text{David}, \text{ULB}) > U(\text{David}, \text{UCL})$$

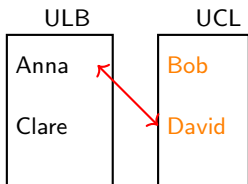
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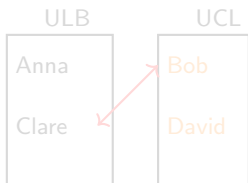
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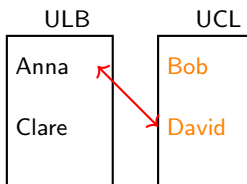
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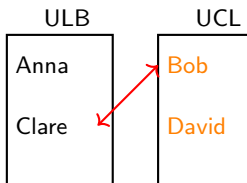
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# Social welfare functions

- Call  $W : \mathcal{M}^* \rightarrow \mathbb{R}$  the social welfare associated to a stable school assignment.
- (Peragine, 2004)

$$W(\mu) = \sum_{i \in I} \omega_i U(i, \mu(i))$$

such that

- ▶ for all  $i, i' \in t$ ,  $\omega_i = \omega_{i'} > 0$ ; and
- ▶ for all  $i \in t$ ,  $i' \in t'$  and  $t, t' \in T$ , if  $t \triangleleft t'$ , then  $\omega_i > \omega_{i'}$ .

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# The SOE algorithm

- The Stable Opportunity Egalitarian (SOE) algorithm maximizes Social Welfare within the set of stable matchings
- Overcomes NP-completeness, efficient running time
- Realizes Affirmative Actions under standard stability notions
- Non technical intuition:
  - ▶ Start from Student optimal Deferred Acceptance
  - ▶ Move to another stable allocation (*reallocate students*) only if it improves Social Welfare

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# Application

# Empirical strategy

- A simulation based on INVALSI data (2021-2022)
- The smallest Italian province: Trieste (212.5 km<sup>2</sup>).
- High school tracks: Scientific (549 students - 4 schools), Humanitarian (161 students - 2 schools), Technical (363 students - 4 schools) and Professional (81 students - 2 schools).
- 5 groups according to Economic, Social and Cultural Status (ESCS)

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# Match quality and preferences

- Match quality: average performance of similar students
- $U(i, s) = \mathbb{E} [\text{Score}_j | \mu(j) = s, j \in t(i)]$
- Simulate (500 times) preferences of students and schools.
- School currently attended is always top ranked in student preference
- We compute

$$W(\mu) = \sum_{i \in I} \omega_i U(i, \mu(i))$$

with  $\omega_i \in \{2, 1.6, 1.4, 1.2, 1\}$

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# Results - Social Welfare

Table: Social welfare difference with respect to Student optimal Deferred Acceptance

Assignment	Scientific	Humanitarian	Technical	Professional
<b>Efficient and Fair</b>	4.76 *** (0 , 35)	2.16 *** (0 , 13)	5.84 *** (0 , 40.4)	2.31 *** (0 , 17.2)

**Description:** The table reports the average difference between social welfare in the considered matching and in Student optimal DA. Minimum and Maximum in parenthesis. P-values: \* ( $p < 0.10$ ), \*\* ( $p < 0.05$ ), \*\*\* ( $p < 0.01$ ).

*Source:* Own elaboration.

## Results - Inequality of opportunity

Table: Inequality of opportunity difference with respect to Student optimal Deferred Acceptance

Assignment	Scientific	Humanitarian	Technical	Professional
<b>Efficient and Fair</b>	-0.02 *** (-0.09 , 0.02)	0 *** (0 , 0.05)	-0.02 *** (-0.11 , 0.1)	0 ** (-0.07 , 0.13)

**Description:** The table reports the average difference between relative IOp in the considered matching and in Student optimal DA. Minimum and Maximum in parenthesis. P-values: \* ( $p < 0.10$ ), \*\* ( $p < 0.05$ ), \*\*\* ( $p < 0.01$ ).

*Source:* Own elaboration.

# Result - Frequency Student Optimal

Table: Probability of Student optimal Deferred Acceptance equal to Assignment.

Assignment	Scientific	Humanitarian	Technical	Professional
<b>Efficient and Fair</b>	0.492	0.422	0.392	0.436

**Description:** The table reports the share of times (out of 500) Student optimal DA has been found to solve the above allocative problems.

*Source:* Own elaboration.

# Conclusion

- Merge Equality of Opportunity and Matching Theory.
- New concept of fairness in school choice settings.
- Define the problem as maximization of a Social Welfare function under stability
- The SOE algorithm solves it efficiently
- Simulation shows potentials for practical implementation
- Many more applications of SOE:
  - ▶ doctor-hospital;
  - ▶ refugees allocation;
  - ▶ equalizing hospital's opportunity for good doctors, to reduce regional disparities in the health care system;
  - ▶ any many-to-one matching problem which can be expressed as maximization of a linear social welfare function over the set of stable matchings

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  - ▶ any many-to-one matching problem which can be expressed as maximization of a linear social welfare function over the set of stable matchings

# Thank you!

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